Smalltalk: The Birthday Problem

Brian Heinold

Mount St. Mary's University

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A Different Problem: How many people have to be in a room with you in order for there to be a 50-50 chance that someone has the same birthday as you?

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• Imagine we have 10 people in a room with different birthdays. Someone new walks into the room. What is the probability their birthday matches someone else's in the room?

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- Etc.
- How many 3% chances can we keep taking before one of them succeeds?
- And with each new person, that 3% gradually increases.

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- In order for there to be no shared birthday, all 253 of those pairs need to work.
- Things just grow from there:
 - With 50 people there are 1225 possible pairs.
 - With 100 people there are 4950 possible pairs

Let's look at a simulation.



Simulation code

```
from random import choice
from time import sleep
months = ["January", "February", "March", "April", "May",
          "June", "July", "August", "September", "October",
          "November". "December"]
days = [31, 28, 31, 30, 31, 30, 31, 31, 30, 31, 30, 31]
days_of_year = []
for i in range(12):
    for j in range(1, days[i]+1):
        days_of_year.append(months[i] + ' ' + str(j)))
prev = set()
c = 0
while True:
    x = choice(days_of_vear)
    print('{:2d}. {:12s} {:3.1f}% chance of repeat on next person'\
          .format(c, x, 100*(c+1)/365))
        if x in prev:
        break
    prev.add(x)
    c += 1
    sleep(.3)
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```

0.	April 11	0.3%	chance	of	repeat	on	next	person
1.	May 8	0.5%	chance	of	repeat	on	next	person
2.	May 1	0.8%	chance	of	repeat	on	next	person
3.	November 16	1.1%	chance	of	repeat	on	next	person
4.	September 17	1.4%	chance	of	repeat	on	next	person
5.	January 30	1.6%	chance	of	repeat	on	next	person
6.	February 2	1.9%	chance	of	repeat	on	next	person
7.	July 8	2.2%	chance	of	repeat	on	next	person
8.	August 17	2.5%	chance	of	repeat	on	next	person
9.	February 8	2.7%	chance	of	repeat	on	next	person
10.	April 11	3.0%	chance	of	repeat	on	next	person

Simulation 2

0.	August 6	0.3%	chance	of	repeat	on	next	person
1.	January 15	0.5%	chance	of	repeat	on	next	person
2.	April 14	0.8%	chance	of	repeat	on	next	person
3.	February 18	1.1%	chance	of	repeat	on	next	person
4.	July 10	1.4%	chance	of	repeat	on	next	person
5.	August 31	1.6%	chance	of	repeat	on	next	person
6.	September 28	1.9%	chance	of	repeat	on	next	person
7.	July 30	2.2%	chance	of	repeat	on	next	person
8.	March 20	2.5%	chance	of	repeat	on	next	person
9.	December 30	2.7%	chance	of	repeat	on	next	person
10.	May 1	3.0%	chance	of	repeat	on	next	person
11.	July 23	3.3%	chance	of	repeat	on	next	person
12.	April 19	3.6%	chance	of	repeat	on	next	person
13.	April 2	3.8%	chance	of	repeat	on	next	person
14.	June 19	4.1%	chance	of	repeat	on	next	person
15.	October 24	4.4%	chance	of	repeat	on	next	person
16.	May 22	4.7%	chance	of	repeat	on	next	person
17.	December 9	4.9%	chance	of	repeat	on	next	person
18.	July 26	5.2%	chance	of	repeat	on	next	person
19.	April 19	5.5%	chance	of	repeat	on	next	person

Simulation 3

0.	June 9	0.3%	chance	of	repeat	on	next	person
1.	January 15	0.5%	chance	of	repeat	on	next	person
2.	March 16	0.8%	chance	of	repeat	on	next	person
3.	April 4	1.1%	chance	of	repeat	on	next	person
4.	November 8	1.4%	chance	of	repeat	on	next	person
5.	September 26	1.6%	chance	of	repeat	on	next	person
6.	February 13	1.9%	chance	of	repeat	on	next	person
7.	January 10	2.2%	chance	of	repeat	on	next	person
8.	October 24	2.5%	chance	of	repeat	on	next	person
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13.	January 26	3.8%	chance	of	repeat	on	next	person
14.	December 13	4.1%	chance	of	repeat	on	next	person
15.	May 11	4.4%	chance	of	repeat	on	next	person

Simulation 4

0.	February 25	0.3% chance of repeat on next person
1.	June 23	0.5% chance of repeat on next person
2.	September 12	0.8% chance of repeat on next person
3.	January 10	1.1% chance of repeat on next person
4.	May 9	1.4% chance of repeat on next person
5.	September 25	1.6% chance of repeat on next person
6.	October 20	1.9% chance of repeat on next person
7.	January 24	2.2% chance of repeat on next person
8.	April 14	2.5% chance of repeat on next person
9.	April 20	2.7% chance of repeat on next person
10.	June 11	3.0% chance of repeat on next person
11.	November 10	3.3% chance of repeat on next person
12.	March 30	3.6% chance of repeat on next person
45.	June 8	12.6% chance of repeat on next person
46.	February 24	12.9% chance of repeat on next person
47.	November 11	13.2% chance of repeat on next person
48.	April 26	13.4% chance of repeat on next person
49.	March 3	13.7% chance of repeat on next person
50.	June 20	14.0% chance of repeat on next person
51.	August 6	14.2% chance of repeat on next person
52.	January 12	14.5% chance of repeat on next person
53.	April 14	14.8% chance of repeat on next perso

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# in room	Probability of two people sharing a birthday
5	2.7%
10	11.7%
20	41.1%
30	70.6%
40	89.1%
50	97.0%
60	99.4%
70	99.9%
80	99.99%
90	99.999%
100	99.99997%

Probability comparison

# in room	Any two	exactly yours			
5	2.7%	1.4%			
10	11.7%	2.7%			
20	41.1%	5.3%			
30	70.6%	7.9%			
40	89.1%	10.4%			
50	97.0%	12.8%			
60	99.4%	15.2%			
70	99.9%	17.5%			
80	99.99%	19.7%			
90	99.999%	21.9%			
100	99.99997%	24.0%			

With k people in a room, the probability p of a match is

$$p = 1 - \frac{364}{365} \cdot \frac{363}{365} \cdot \frac{362}{365} \cdot \dots \cdot \frac{365 - (k-1)}{365}.$$

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Use the multiplication rule to get that probability.

• With *k* people in a room, the probability *p* of no shared birthdays is

$$\begin{split} p = & 1 - \frac{364}{365} \cdot \frac{363}{365} \cdot \frac{362}{365} \cdot \dots \cdot \frac{365 - (k-1)}{365} \\ = & 1 - \left(1 - \frac{1}{365}\right) \left(1 - \frac{2}{365}\right) \left(1 - \frac{3}{365}\right) \dots \left(1 - \frac{k-1}{365}\right) \\ \approx & 1 - e^{-1/365} e^{-2/365} \dots e^{-(k-1)/365} \\ = & 1 - e^{-k(k-1)/(2 \cdot 365)} \\ \approx & 1 - e^{-k^2/(2 \cdot 365)} \end{split}$$

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• Invert this to get the number of people needed for there to be a probability *p* of a repeat:

$$k \approx \sqrt{2 \cdot 365 \ln\left(\frac{1}{1-p}\right)}$$

Generalizing the birthday problem

There's nothing special about birthdays.

If we generate randomly from a set of n items, the probability of a repeat after k items is

$$p = 1 - \frac{n-1}{n} \cdot \frac{n-2}{n} \cdot \frac{n-3}{n} \cdot \dots \cdot \frac{n-(k-1)}{n}$$

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$$k = \sqrt{2 \cdot n \ln\left(\frac{1}{1-p}\right)}$$

And, the most commonly used rule of thumb is that after \sqrt{n} things are generated, repeats are fairly likely.

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$$\sqrt{2 \cdot 1000 \ln\left(\frac{1}{1 - .25}\right)} \approx 24$$

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Solution About how many before a repeat is likely?
An example

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 About how many before a repeat is likely? Quick estimate: √1000 ≈ 32

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$$p = 1 - \frac{4999}{5000} \cdot \frac{4998}{5000} \cdot \frac{4997}{5000} \cdot \dots \cdot \frac{4951}{5000} \approx 22\%$$

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• Quick estimate: repeats are likely after around $\sqrt{5000} \approx 70$ songs

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- But how likely is it that over the course of your life, you will be dealt a hand that you had been dealt at some time in the past?
- That's the birthday problem.
- After roughly $\sqrt{2,600,000} \approx 1600$ hands, repeats are likely.

The quick estimate $\sqrt{2,600,000} \approx 1600$ is nice because it gives an order of magnitude for when we should expect repeats. For example:

Probability	# of hands for a repeat
5%	517
20%	1078
39%	1603
50%	1899
75%	2685
99%	4894
99.999%	7738

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- But the birthday problem matters. After only $\sqrt{18,000,000,000,000} \approx 4$ billion transactions, a repeat is likely.
- If this is a large internet site, 4 billion transactions is quite possible.

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- Basic idea is that you feed them a string and they return a fixed length output.
- MD5 is one well-used hash function that returns 64-bit outputs. Example hashes:

"smalltalk" 90945bf1d2c52618e38eada42b86086b

Chapter 1 of A Tale of Two Cities ff5eb755a31ea99fca07b9fba8dd1d07

Chapter 1 of *A Tale of Two Cities* with first letter changed to *Z* and everything else left intact a5d7c988d5d87d1b28e1d85e77110cd1

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- For these reasons, hash functions are used as fingerprints.
- For instance, the hash of Chapter 1 of *A Tale of Two Cities*, ff5eb755a31ea99fca07b9fba8dd1d07, is highly unlikely to be the output of any other string created by humans in history.
- Why? A 64-bit hash function has 2⁶⁴ ≈ 18,000,000,000,000,000 possible outputs, far more than the number of strings created by humans in history.

• Hashes are used in digital signatures: A person sending you a document can compute its hash, and you can compute its hash when you get the document. If the hashes match, then you know the document you got is the same as what the person sent.

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- Birthday problem: If we generate $\sqrt{18 \text{ quintillion}} \approx 4 \text{ billion}$ strings, it is likely that some pair of them have the same MD5 hash.

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 - I generate a few billion subtle variations on the \$10 contract as well.
- Eventually, I will get a hash collision where one of the \$100 variations matches one of the \$10 variations.
- Since they both have the same hash, you might think you have the real one, but you actually have a fake one.

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- This is because $\sqrt{2^{128}} = 2^{64} = 10^{18}$, which is a lot of variations to make, but within the reach of well-funded nation states.
- So 256-bit hash functions are recommended.

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- If you do, it's trivially easy to crack.

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- This would be fine, except they used 24-bit IVs.
- A 24-bit IV would mean 16 million possible encryption keys
- Now the birthday problem comes in: $\sqrt{2^{24}} = 2^{12} = 4096$
- After 4096 messages, repeated keys are likely.
- This is the amount of traffic you get on a typical network in a few minutes.

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- That machine asks other machines on the internet for the answer.
- If an attacker can return an answer before the real answer arrives, then the attacker's answer will be accepted.
- The attacker can use this to make it so that when you go to gmail.com, you are actually directed to a fake gmail site, where they can phish your username and password.

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- There are $2^{16} = 65536$ possible IDs, so guessing is hard.
- One more thing: the machine your ISP uses stores the answers to the DNS queries it does, so that if someone else with the same provider as you recently went to msmary.edu, the DNS machine will store the resulting IP address to save time from having to ask remote machines for the answer.

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- Most of those won't have matching IDs, but the attacker just needs one fake answer to match one fake request and then that result will be stored in memory for a while, affecting many of the service provider's customers.
- By the birthday problem, they need only create around $\sqrt{2^{16}} = 2^8 = 256$ bogus requests, which is easy.

A fun little math problem turns out to have big consequences in computer security. The examples shown here are just a few of many more.

Thanks for your attention!